Reasoning about the Elementary Functions of Complex Analysis*

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Abstract. There are many problems with the simplification of elementary functions, particularly over the complex plane. Systems tend to make major errors, or not to simplify enough. In this paper we outline the "unwinding number" approach to such problems, and show how it can be used to prevent errors and to systematise such simplification, even though we have not yet reduced the simplification process to a complete algorithm. The unsolved problems are probably more amenable to the techniques of artificial intelligence and theorem proving than the original problem of complex-variable analysis.
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1 Introduction

The elementary functions are traditionally thought of as log, exp and the trigonometric and hyperbolic functions (and their inverses). This list should include powering (to non-integral powers) and also the n-th root. These functions are built in, to a greater or lesser extent, to many computer algebra systems (not to mention other programming languages [8, 12]), and are heavily used. However, reasoning with them is more difficult than is usually acknowledged, and all algebra systems have one, sometimes both, of the following defects:

– they make mistakes, be it the traditional schoolchild one

$$1 = \sqrt{1} = \sqrt{(-1)^2} = -1 \quad (1)$$

or more subtle ones (see footnote 6);

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– they fail to perform obvious simplifications, leaving the user with an impossible mess when there "ought" to be a simpler answer. In fact, there are two possibilities here: maybe there is a simpler equivalent that the system has failed to find, but maybe there isn't, and the simplification that the user wants is not actually valid, or is only valid outside an exceptional set. In general, the user is not informed what the simplification might have been, nor what the exceptional set is.

Faced with these problems, the user of the algebra system is not convinced that the result is correct, or that the algebra system in use understands the functions with which it is reasoning. An ideal algebra system would never generate incorrect results, and would simplify the results as much as practicable, even though perfect simplification is impossible, and not even totally well-defined: is $1 + x + \cdots + x^{1000}$ "simpler" than $(x^{1001} - 1)/(x - 1)$?

Throughout this paper, z and its decorations indicate a complex variable, while x, y and t indicate real variables. The symbol $\Im$ denotes the imaginary part, and $\Re$ the real part, of a complex number. For the purposes of this paper, the precise definitions of the inverse elementary functions in terms of log are those of [4]: these are reproduced in Appendix A for ease of reference.

2 The Problem

The fundamental problem is that log is multi-valued: since $\exp(2\pi i) = 1$, its inverse is only valid up to adding any multiple of $2\pi i$. This ambiguity is traditionally resolved by making a *branch cut*: usually [1, p. 67] the branch cut $(-\infty, 0]$, and the rule (4.1.2) that

$$-\pi < \Im \log z \leq \pi. \tag{2}$$

This then completely specifies the behaviour of log: on the branch cut it is continuous with the positive imaginary side of the cut, i.e. counter-clockwise continuous in the sense of [10].

What are the consequences of this definition[1]? From the existence of branch cuts, we get the problem of a lack of continuity:

$$\lim_{y \to 0^-} \log(x + iy) \neq \log x : \tag{3}$$

for $x < 0$ the limit is $\log x - 2\pi i$. Related to this is the fact that

$$\log \overline{z} \neq \overline{\log z} \tag{4}$$

[1] Which we do not contest: it seems that few people today would support the rule one of us (JHD) was taught, viz. that $0 \leq \Im \log z < 2\pi$. The placement of the branch cut is "merely" a notational convention, but an important one. If we wanted a function that behaves like log but with this cut, we could consider $\underbrace{\log}_{[0,2\pi)}(z) = \log(-1) - \log(-1/z)$ instead. We note that, until 1925, astronomers placed the branch cut between one day and the next at noon [7, vol. 15 p. 417].

on the branch cut: instead $\log \overline{z} = \overline{\log z} + 2\pi i$ on the cut. Similarly,

$$\log\left(\frac{1}{z}\right) \neq -\log z \tag{5}$$

on the branch cut: instead $\log(1/z) = -\log z + 2\pi i$ on the cut.

Although not normally explained this way, the problem with (1) is a consequence of the multi-valued nature of log: if we define (as for the purposes of this paper we do)

$$\sqrt{z} = \exp\left(\frac{1}{2}\log z\right), \tag{6}$$

then $-\pi/2 < \Im\sqrt{z} \leq \pi/2$. On the real line, this leads to the traditional resolution of (1), namely that $\sqrt{x^2} = |x|$.

Three families of solutions have been proposed to these problems.

- Prof. W. Kahan points out that the concept of a "signed zero"[2] [9] (for clarity, we write the positive zero as 0^+ and the negative one as 0^-) can be used to solve the above problems, if we say that, for $x < 0$, $\log(x + 0^+ i) = \log x + \pi i$ whereas $\log(x + 0^- i) = \log x - \pi i$. Equation (3) then becomes an equality for all x, interpreting the x on the right as $x + 0^- i$. Similarly, (4) and (5) become equalities throughout. Attractive though this proposal is, it does not answer the fundamental question as far as the designer of a computer algebra system is concerned: what to do if the user types $\log(-1)$.
- The authors of [5] point out that most "equalities" do not hold for the complex logarithm, e.g. $\log(z^2) \neq 2\log z$ (try $z = -1$), and its generalisation

  $$\log(z_1 z_2) \neq \log z_1 + \log z_2. \tag{7}$$

  The most fundamental of all non-equalities is $z \stackrel{?}{=} \log \exp z$, whose most obvious violation is at $z = 2\pi i$. (A similar point was made in [2], where the correction term is called the "adjustment".) They therefore propose to formalise the violation of this equality by introducing the *unwinding number* $\mathcal{K}$, defined[3] by

  $$\mathcal{K}(z) = \frac{z - \log e^z}{2\pi i} = \left\lceil \frac{\Im z - \pi}{2\pi} \right\rceil \in \mathbf{Z} \tag{8}$$

  (note that the apparently equivalent definition $\left\lfloor \frac{\Im z + \pi}{2\pi} \right\rfloor$ differs precisely on the branch cut for log as applied to $\exp z$).

[2] One could ask why zero should be special and have two values. The answer seems to be that all the branch cuts we need to consider are on either the real or imaginary axes, so the side to which the branch cut adheres depends on the sign of the imaginary or real part, including the sign of zero. To handle other points similarly would require the arithmetic of non-standard analysis.

[3] Note that the sign convention here is the opposite to that of [5], which defined $\mathcal{K}(z)$ as $\left\lfloor \frac{\pi - \Im z}{2\pi} \right\rfloor$: the authors of [5] recanted later to keep the number of -1s occurring in formulae to a minimum. We could also change "unwinding" to "winding" when we make that sign change; but "winding number" is in wide use for other contexts, and it seems best to keep the existing terminology.

This definition has several attractive features: $\mathcal{K}(z)$ is integer-valued, and familiar in the sense that "everyone knows" that the multivalued logarithm can be written as the principal branch "plus $2\pi ik$ for some integer k"; it is single-valued; and it can be computed by a formula not involving logarithms. It does have a numerical difficulty, namely that you must decide if the imaginary part is an odd integer multiple of π or not, and this can be hard (or impossible in some exact arithmetic contexts), but the difficulty is inherent in the problem and cannot be repaired e.g. by putting the branch cuts elsewhere.
Some correct identities for elementary functions using $\mathcal{K}$ are given in Table 1.

1. $z = \log e^z + 2\pi i\mathcal{K}(z)$.
2. $\mathcal{K}(a \log z) = 0\,\forall z \in \mathbf{C}$ if and only if $-1 < a \leq 1$.
3. $\log z_1 + \log z_2 = \log(z_1 z_2) + 2\pi i\mathcal{K}(\log z_1 + \log z_2)$.
4. $a \log z = \log z^a + 2\pi i\mathcal{K}(a \log z)$.
5. $z^{ab} = (z^a)^b e^{2\pi i b\mathcal{K}(a \log z)}$.

Table 1. Some correct identities for logarithms and powers using $\mathcal{K}$.

(7) can then be rescued as

$$\log(z_1 z_2) = \log z_1 + \log z_2 - 2\pi i\mathcal{K}(\log z_1 + \log z_2). \tag{9}$$

Similarly (4) can be rescued as

$$\log \overline{z} = \overline{\log z} - 2\pi i\mathcal{K}\left(\overline{\log z}\right). \tag{10}$$

Note that, as part of the algebra of $\mathcal{K}$, $\mathcal{K}(\overline{\log z}) = \mathcal{K}(-\log z) \neq \mathcal{K}(\log 1/z)$. $\mathcal{K}(z)$ depends only on the imaginary part of z.

– Although not formally proposed in the same way in the computational community, one possible solution, often found in texts in complex analysis, is to accept the multi-valued nature of these functions (we adopt the common convention of using capital letters, e.g. Ln, to denote the multi-valued function), defining, for example

$$\mathrm{Arcsin}\, z = \{y \mid \sin y = z\}.$$

This leads to $\sqrt{z^2} = \{\pm z\}$, which has the advantage that it is valid throughout $\mathbf{C}$. Equation 7 is then rewritten as

$$\mathrm{Ln}(z_1 z_2) = \mathrm{Ln}\, z_1 + \mathrm{Ln}\, z_2\,, \tag{11}$$

where addition is addition of sets ($A + B = \{a + b : a \in A, b \in B\}$) and equality is set equality[4].

[4] "The equation merely states that the sum of one of the (infinitely many) logarithms of z_1 and one of the (infinitely many) logarithms of z_2 can be found among the

However, it seems to lead in practice to very large and confusing formulae. More fundamentally, this approach does not say what will happen when the multi-valued functions are replaced by the single-valued ones of numerical programming languages.

A further problem that has not been stressed in the past is that this approach suffers from the same aliasing problem that naïve interval arithmetic does [6]. For example,

$$\mathrm{Ln}(z^2) = \mathrm{Ln}\, z + \mathrm{Ln}\, z \neq 2\,\mathrm{Ln}\, z\ ,$$

since $2\,\mathrm{Ln}(z) = \{2\log(z) + 4k\pi i : k \in \mathbf{Z}\}$, but $\mathrm{Ln}(z) + \mathrm{Ln}(z) = \{2\log(z) + 2k\pi i : k \in \mathbf{Z}\}$: indeed if $z = -1$, $\log(z^2) \notin 2\,\mathrm{Ln}(z)$. Hence this method is unduly pessimistic: it may fail to prove some identities that are true.

3 The rôle of the Unwinding Number

We claim that the unwinding number provides a convenient formalism for reasoning about these problems. Inserting the unwinding number systematically allows one to make "simplifying" transformations that *are* mathematically valid. The unwinding number can be evaluated at any point, either symbolically or via guaranteed arithmetic: since we know it is an integer, in practice little accuracy is necessary. Conversely, removing unwinding numbers lets us genuinely "simplify" a result. We describe insertion and removal as separate steps, but in practice every unwinding number, once inserted by a "simplification" rule, should be eliminated as soon as possible. We have thus defined a concrete goal for mathematically valid simplification.[5]

The following section gives examples of reasoning with unwinding numbers. Having motivated the use of unwinding numbers, the subsequent sections deal with their insertion (to preserve correctness) and their elimination (to simplify results).

4 Examples of Unwinding Numbers

This section gives certain examples of the use of unwinding numbers. We should emphasise our view that an ideal computer algebra system should do this manipulation for the user: certainly inserting the unwinding numbers where necessary, and preferably also removing/simplifying them where it can.

4.1 Forms of arccos

The following example is taken from [4], showing that two alternative definitions of arccos are in fact equal:

(infinitely many) logarithms of $z_1 z_2$, and conversely every logarithm of $z_1 z_2$ can be represented as a sum of this kind (with a suitable choice of [elements of] $\mathrm{Ln}\, z_1$ and $\mathrm{Ln}\, z_2$)." [3, pp. 259–260] (our notation).

[5] Just to remove the terms with unwinding numbers, as is done in some software systems, could be called "over-simplification."

Theorem 1.

$$\frac{2}{i}\log\left(\sqrt{\frac{1+z}{2}}+i\sqrt{\frac{1-z}{2}}\right) = -i\log\left(z+i\sqrt{1-z^2}\right). \tag{12}$$

First we prove the correct (and therefore containing unwinding numbers) version of $\sqrt{z_1 z_2} \stackrel{?}{=} \sqrt{z_1}\sqrt{z_2}$.

Lemma 1.

$$\sqrt{z_1 z_2} = \sqrt{z_1}\sqrt{z_2}(-1)^{\mathcal{K}(\log z_1+\log z_2)}. \tag{13}$$

Proof.

$$\begin{aligned}\sqrt{z_1 z_2} &= \exp\left(\frac{1}{2}(\log(z_1 z_2))\right)\\ &= \exp\left(\frac{1}{2}(\log z_1+\log z_2-2\pi i\mathcal{K}(\log z_1+\log z_2))\right)\\ &= \sqrt{z_1}\sqrt{z_2}\exp(-\pi i\mathcal{K}(\log z_1+\log z_2))\\ &= \sqrt{z_1}\sqrt{z_2}(-1)^{\mathcal{K}(\log z_1+\log z_2)}\end{aligned}$$

Lemma 2. *Whatever the value of z,*

$$\sqrt{1-z}\sqrt{1+z} = \sqrt{1-z^2}.$$

This is a classic example of a result that is "obvious": the schoolchild just squares both sides, but in fact that loses information, and the identity requires proof. To show this, consider the apparently similar "result"[6]:

$$\sqrt{-i-z}\sqrt{-i+z} \stackrel{?}{=} \sqrt{-1-z^2}.$$

If we take $z = i/2$, the left-hand side becomes $\sqrt{-3i/2}\,\sqrt{-i/2}$: the inputs to the square roots[7] have $\arg = -\pi/2$, so the square roots themselves have $\arg = -\pi/4$, and the product has $\arg = -\pi/2$, and therefore is $-i\sqrt{3}/2$. The right-hand side is $\sqrt{-3/4} = i\sqrt{3}/2$.

Proof. It is sufficient to show that the unwinding number term in lemma 1 is zero. Whatever the value of z, $1+z$ and $1-z$ have imaginary parts of opposite signs. Without loss of generality, assume $\Im z \geq 0$. Then $0 \leq \arg(1+z) \leq \pi$ and $-\pi < \arg(1-z) \leq 0$. Therefore their sum, which is the imaginary part of $\log(1+z)+\log(1-z)$, is in $(-\pi,\pi]$. Hence the unwinding number is indeed zero.

Proof of Theorem 1. Now

$$\left(\sqrt{\frac{1+z}{2}}+i\sqrt{\frac{1-z}{2}}\right)^2 = z+i\sqrt{1-z}\sqrt{1+z} = z+i\sqrt{1-z^2}$$

[6] Maple V.5, in the absence of an explicit declaration that z is complex, will say that the two are almost never equal, with the difference being $-2i\sqrt{1-z^2}$, but in fact at $z = 2i$, the two are equal.

[7] One is tempted to say "arguments of the square root", but this is easily confused with the function arg; we use 'inputs' instead.

by the previous lemma. Also $2\log a=\log(a^2)$ if $\mathcal{K}(2\log a)=0$, so we need only show this last stipulation, i.e. that

$$-\frac{\pi}{2}<\arg\left(\sqrt{\frac{1+z}{2}}+i\sqrt{\frac{1-z}{2}}\right)\leq\frac{\pi}{2}.$$

This is trivially true at $z=0$. If it is false at any point, say z_0, then a path from z_0 to 0 must pass through a z where $\left|\arg\left(\sqrt{(1+z)/2}+i\sqrt{(1-z)/2}\right)\right|=\pi/2$, i.e. $\sqrt{(1+z)/2}+i\sqrt{(1-z)/2}=it$ for $t\in\mathbf{R}$, because, first, arg is continuous for $|z|\leq\pi/2$, and indeed for $|z|<\pi$, and, second, that the inputs to arg are themselves discontinuous only on $z>1$ and $z<-1$, and on these half-lines, the arguments in question are 0 and $\pi/2$, which are acceptable. Coming back to the continuity along the path, we find that by squaring both sides, $z+i\sqrt{1-z^2}=-t^2$, i.e. $(z+t^2)^2=-(1-z^2)$. Hence $2zt^2+t^4=-1$, so $z=-(1+t^4)/(2t^2)\leq-1$, and in particular is real. On this half-line, as stated before, the argument in question is $+\pi/2$, which is acceptable. Hence the argument never leaves the desired range, and the theorem is proved.

4.2 arccos and arccosh

$\cos(z)=\cosh(iz)$, so we can ask whether the corresponding relation for the inverse functions, $\operatorname{arccosh}(z)=i\arccos(z)$, holds. This is known in [4] as the "couthness" of the arccos/arccosh definitions. The problem reduces, using equations (20) and (26), to

$$2\log\left(\sqrt{\frac{z+1}{2}}+\sqrt{\frac{z-1}{2}}\right)\stackrel{?}{=}i\left(\frac{2}{i}\log\left(\sqrt{\frac{1+z}{2}}+i\sqrt{\frac{1-z}{2}}\right)\right),$$

i.e.

$$\log\left(\sqrt{\frac{z+1}{2}}+\sqrt{\frac{z-1}{2}}\right)\stackrel{?}{=}\log\left(\sqrt{\frac{1+z}{2}}+i\sqrt{\frac{1-z}{2}}\right).$$

Since $\log a=\log b$ implies $a=b$ (n.b. this is not true for exp, which is part of the point of this paper), this reduces to

$$\sqrt{\frac{z-1}{2}}\stackrel{?}{=}i\sqrt{\frac{1-z}{2}}=\sqrt{-1}\sqrt{\frac{1-z}{2}}.$$

By lemma 1, the right-hand side reduces to $\sqrt{\frac{z-1}{2}}(-1)^{-\mathcal{K}(\log(-1)+\log(\frac{z-1}{2}))}$. Hence the two are equal if, and only if, the unwinding number is even (and therefore zero). This will happen if, and only if, $\arg\left(\frac{z-1}{2}\right)\leq 0$, i.e. $\Im z<0$ or $\Im z=0$ and $z>1$.

4.3 arcsin and arctan

The aim of this section is to prove the correct expression for arcsin in terms of arctan. We note that we need to add unwinding number terms to deal with the two cuts $\Re z < -1$, $\Im z = 0$ and $\Re z > 1$, $\Im z = 0$.

Theorem 2.

$$\arcsin z = \arctan \frac{z}{\sqrt{1-z^2}} + \pi\mathcal{K}(-\log(1+z)) - \pi\mathcal{K}(-\log(1-z)). \tag{14}$$

We start from equations (19) and (21). Then

$$\begin{aligned}
2i\ \arctan \frac{z}{\sqrt{1-z^2}} &= \log\left(1 + i\frac{z}{\sqrt{1-z^2}}\right) - \log\left(1 - i\frac{z}{\sqrt{1-z^2}}\right) \\
&= \log\left([1 + i\frac{z}{\sqrt{1-z^2}}]/[1 - i\frac{z}{\sqrt{1-z^2}}]\right) \\
&\quad +2\pi i\mathcal{K}\left(\log(1 + i\frac{z}{\sqrt{1-z^2}}) - \log(1 - i\frac{z}{\sqrt{1-z^2}})\right) \\
&= \log[iz + \sqrt{1-z^2}]^2 \\
&\quad +2\pi i\mathcal{K}\left(\log(1 + i\frac{z}{\sqrt{1-z^2}}) - \log(1 - i\frac{z}{\sqrt{1-z^2}})\right) \\
&= 2i\arcsin(z) \\
&\quad -2\pi i\mathcal{K}\left(2\log(iz + \sqrt{1-z^2})\right) \\
&\quad +2\pi i\mathcal{K}\left(\log(1 + i\frac{z}{\sqrt{1-z^2}}) - \log(1 - i\frac{z}{\sqrt{1-z^2}})\right)
\end{aligned}$$

The tendency for $\mathcal{K}$ factors to proliferate is clear. To simplify we proceed as follows. Consider first the term

$$\mathcal{K}\left(2\log(iz + \sqrt{1-z^2})\right) .$$

For $|z| < 1$, the real part of the input to the logarithm is positive and hence has argument in $(-\pi/2, \pi/2)$; therefore $\mathcal{K} = 0$. For $|z| > 1$, we solve for the critical case in which the input to $\mathcal{K}$ is $-i\pi$ and find only $z = r\exp(i\pi)$, with $r > 1$. Therefore

$$\mathcal{K}(2\log(iz + \sqrt{1-z^2})) = \mathcal{K}(-\log(1+z)) .$$

Repeating the procedure with

$$\mathcal{K}\left(\log(1 + iz/\sqrt{1-z^2}) - \log(1 - iz/\sqrt{1-z^2})\right)$$

shows that $\mathcal{K} \neq 0$ only for $z > 1$. Therefore

$$\mathcal{K}\left(\log(1 + iz/\sqrt{1-z^2}) - \log(1 - iz/\sqrt{1-z^2})\right) = \mathcal{K}(-\log(1-z))$$

and so finally we get

$$\arctan \frac{z}{\sqrt{1-z^2}} = \arcsin(z) - \pi\mathcal{K}(-\log(1+z)) + \pi\mathcal{K}(-\log(1-z)) , \tag{15}$$

and this cannot be simplified further.

5 The Unwinding Number: Insertion

We have seen that the systematic insertion of unwinding numbers while applying many "simplification" rules is necessary for mathematical correctness.

Unwinding numbers are normally inserted by use of equation (9) and its converse:

$$\log\left(\frac{z_1}{z_2}\right) = \log z_1 - \log z_2 - 2\pi\mathcal{K}\left(\log z_1 - \log z_2\right). \tag{16}$$

Equation (10) may also be used, as may its close relative (also a special case of (16))

$$\log\left(\frac{1}{z}\right) = -\log z - 2\pi\mathcal{K}\left(-\log z\right). \tag{17}$$

In practice, results such as lemma 1 would also be built in to a simplifier.

The definition of $\mathcal{K}$ gives us

$$\log(e^z) = z - 2\pi i\mathcal{K}(z), \tag{18}$$

which is another mechanism for inserting unwinding numbers while "simplifying". The formulae for other inverse functions are given in appendix B.

Many other "identities" among inverse functions require unwinding numbers. For example,

$$\arctan x + \arctan y = \arctan\left(\frac{x+y}{1-xy}\right) + \pi\mathcal{K}\left(2i(\arctan x + \arctan y)\right).$$

6 The Unwinding Number: Removal

It is clearly easier to insert unwinding numbers than to remove them. There are various possibilities for the values of unwinding numbers.

- An unwinding number may be identically zero. This is the case in lemma 2 and theorem 1. The aim is then to prove this.
- An unwinding number may be zero everywhere except on certain branch cuts in the complex plane. This is the case in equation (10), and its relative $\log(1/z) = -\log z - 2\pi i\mathcal{K}(-\log z)$. A less trivial case of this can be seen in equation (14). Derive has a different definition of arctan to eliminate this, so that, for Derive, $\arcsin(z) = \underbrace{\arctan}_{\text{Derive}} \frac{z}{\sqrt{1-z^2}}$. This definition can be related to ours either via unwinding numbers or via $\underbrace{\arctan}_{\text{Derive}}(z) = \overline{\arctan \overline{z}}$. It is often possible to disguise this sort of unwinding number, which is often of the form $\mathcal{K}(-\log(\ldots))$ or $\mathcal{K}(\overline{\log z})$, by resorting to such a "double conjugate" expression, though as yet we have no algorithm for this. Equally, we have no algorithm as yet for the sort of simplification we see in section 4.3.

– An unwinding number may divide the complex plane into two regions, one where it is non-zero and one where it is zero. A typical case of this is given in section 4.2. Here the proof methodology consists in examining the critical case, i.e. when the input to $\mathcal{K}$ has imaginary part $\pm\pi$, and examining when the functions contained in the input to $\mathcal{K}$ themselves have discontinuities.
– An unwinding number may correspond to the usual $+n\pi$: $n \in \mathbf{Z}$ of many trigonometric identities: examples of this are given in appendix B.

7 Conclusion

Unwinding number insertion permits the manipulation of logarithms, square roots etc., as well as the cancellation of functions and their inverses, while retaining mathematical correctness. This can be done completely algorithmically, and we claim this is one way, the only way we have seen, of guaranteeing mathematical correctness while "simplifying".

Unwinding number removal, where it is possible, then simplifies these results to the expected form. This is not a process that can currently be done algorithmically, but it is much better suited to current artificial intelligence techniques than the general problems of complex analysis.

When the unwinding numbers cannot be eliminated, they can often be converted into a case analysis that, while not ideal, is at least comprehensible while being mathematically correct.

More generally, we have reduced the analytic difficulties of simplifying these functions to more algebraic ones, in areas where we hope that artificial intelligence and theorem proving stand a better chance of contributing to the problem.

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A Definition of the Elementary Inverse Functions

These definitions are taken from [4]. They agree with [1, ninth printing], but are more precise on the branch cuts, and agree with Maple with the exception of arccot, for the reasons explained in [4].

$$\arcsin z = -i \log\left(\sqrt{1-z^2} + iz\right). \tag{19}$$

$$\arccos(z) = \frac{\pi}{2} - \arcsin(z) = \frac{2}{i} \log\left(\sqrt{\frac{1+z}{2}} + i\sqrt{\frac{1-z}{2}}\right). \tag{20}$$

$$\arctan(z) = \frac{1}{2i}\left(\log(1+iz) - \log(1-iz)\right). \tag{21}$$

$$\operatorname{arccot} z = \frac{1}{2i} \log\left(\frac{z+i}{z-i}\right) = \arctan\left(\frac{1}{z}\right). \tag{22}$$

$$\operatorname{arcsec}(z) = \arccos(1/z) = -i\log(1/z + i\sqrt{1-1/z^2}), \tag{23}$$

with $\operatorname{arcsec}(0) = \frac{\pi}{2}$.

$$\operatorname{arccsc}(z) = \arcsin(1/z) = -i\log(i/z + \sqrt{1-1/z^2}), \tag{24}$$

with $\operatorname{arccsc}(0) = 0$.

$$\operatorname{arcsinh}(z) = \log\left(z + \sqrt{1+z^2}\right). \tag{25}$$

$$\operatorname{arccosh}(z) = 2\log\left(\sqrt{\frac{z+1}{2}} + \sqrt{\frac{z-1}{2}}\right). \tag{26}$$

$$\operatorname{arctanh}(z) = \frac{1}{2}\left(\log(1+z) - \log(1-z)\right). \tag{27}$$

$$\operatorname{arccoth}(z) = \frac{1}{2}\left(\log(-1-z) - \log(1-z)\right). \tag{28}$$

$$\operatorname{arcsech}(z) = 2\log\left(\sqrt{\frac{z+1}{2z}} + \sqrt{\frac{1-z}{2z}}\right). \tag{29}$$

$$\operatorname{arccsch}(z) = \log\left(\frac{1}{z} + \sqrt{1 + \left(\frac{1}{z}\right)^2}\right), \tag{30}$$

B Formulae for inverse functions

These formulae are taken from [11]. They make use of the secondary function csgn, which we define below in terms of $\mathcal{K}$ and was first defined by Dr. D. E. G. Hare as the piecewise function on the right hand side[8]:

$$\operatorname{csgn}(z) = (-1)^{\mathcal{K}(2\log(z))} = \begin{cases} +1 & \Re(z) > 0 \text{ or } \Re(z) = 0;\ \Im(z) \geq 0 \\ -1 & \Re(z) < 0 \text{ or } \Re(z) = 0;\ \Im(z) < 0 \end{cases}.$$

$$\arcsin(\sin(z)) = \begin{cases} z - 2\pi\mathcal{K}(zi) & \operatorname{csgn}(\cos z) = 1 \\ \pi - z - 2\pi\mathcal{K}(i(\pi - z)) & \operatorname{csgn}(\cos z) = -1 \end{cases}. \tag{31}$$

$$\arccos(\cos z) = \begin{cases} z - 2\pi\mathcal{K}(zi) & \operatorname{csgn}(\sin z) = 1 \\ -z - 2\pi\mathcal{K}(-zi) & \operatorname{csgn}(\sin z) = -1 \end{cases}. \tag{32}$$

$$\arctan(\tan z) = z + \pi\left(\mathcal{K}(-zi - \log\cos z) - \mathcal{K}(zi - \log\cos z)\right) \tag{33}$$

provided $z \neq \frac{\pi}{2} + n\pi$: $n \in \mathbf{Z}$.

$$\operatorname{arcsinh}(\sinh(z)) = \begin{cases} z - 2\pi i\mathcal{K}(z) & \operatorname{csgn}(\cosh z) = 1 \\ i\pi - z - 2\pi i\mathcal{K}(i\pi - z) & \operatorname{csgn}(\cosh z) = -1 \end{cases}. \tag{34}$$

$$\operatorname{arccosh}(\cosh z) = \begin{cases} z - 2\pi\mathcal{K}(z) & \operatorname{csgn}(\sinh z)\cos(n\pi) = 1 \\ -z - 2\pi i\mathcal{K}(-z) & \operatorname{csgn}(\sinh z)\cos(n\pi) = -1 \end{cases} \tag{35}$$

where $n = \mathcal{K}\left(\log(\cosh(z) - 1) + \log(\cosh(z) + 1)\right)$.

$$\operatorname{arctanh}(\tanh z) = z + i\pi\left(\mathcal{K}(z - \log\cosh z) - \mathcal{K}(z - \log\cosh z)\right) \tag{36}$$

provided $z \neq \frac{\pi}{2}i + in\pi$: $n \in \mathbf{Z}$.

[8] This function simplifies $\sqrt{z^2}$ to $z\operatorname{csgn}(z)$. Dr. J. Carette observed that if we put $\omega = \exp(2\pi i/n)$, then the function defined by $\omega^{\mathcal{K}(n\log z)}$ and sometimes abbreviated by $C_n(z)$, that generalizes csgn, is useful in simplifying $(z^n)^{1/n}$ (private communication).